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# A HYBRID MODIFIED GENETIC-NELDER MEAD SIMPLEX ALGORITHM FOR LARGE-SCALE TRUSS OPTIMIZATION

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# ABSTRACT

In this paper a hybrid algorithm based on exploration power of the Genetic algorithms and exploitation capability of Nelder Mead simplex is presented for global optimization of multivariable functions. Some modifications are imposed on genetic algorithm to improve its capability and efficiency while being hybridized with Simplex method. Benchmark test examples of structural optimization with a large number of variables and constraints are chosen to show the robustness of the algorithm.

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## **1. INTRODUCTION**

Recently global optimization has been a concern of researchers, especially when the fitness function is dependent on a large number of variables or it is strictly confined to some constraints. Hybridizing is the only motivation that brings back the hope to reach this goal as long as using only one search algorithm leads to dim results. A combination of two algorithms, in which one explores a promising area likely to contain global minima, and the other exploits the area to find the desired point would be promising if properly performed. Global methods like simulated annealing, tabu search, genetic algorithm, etc are known as efficient search engines to find and localize areas containing global minima, but they are

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very much time consuming while converging to a specified point (if not get trapped in local minimum). On the other hand local methods including Nelder-Mead Simplex, hill climbing etc, are good for exploitation of the search domain. In this study we present a hybrid algorithm based on exploitation (diversification) power of genetic algorithm and exploitation (intensification) feature of the Nelder-Mead Simplex.

The advantage of the simplex search method is that it is straightforward in an algorithmic sense and computationally efficient. However, as a result of using only local information, when they converge to a stationary point, there is no guarantee that the global optimum is found unless the domain in which the global minimum lies is provided. In contrast, a GA method explores the global search space without using local information of promising search directions. But its computational cost is comparatively high. A hybrid algorithm using both of these characteristics (exploration and exploitation) would be promising to find global minima in a broad types of problems. It is practically important to note that each of these algorithms and the new hybrid algorithm presented in this study do not require gradient or Hessian matrix calculations, and therefore it does not suffer from the weakness of classical optimization methods. The main goals of the present hybrid algorithm are as follows:

- Reliability: A proper functioning of exploration and exploitation of the search domain (sometimes referred to as "diversification" and "intensification") in order to find the true global minima.
- Efficiency: Using simple but effective combination to reduce the total amount of function evaluation.

Up to now we introduced the concepts related to this subject. Such general aspects can also be found in Refs. [1-6].

Proper numerical example plays an important role to reveal the robustness of a new algorithm. The large number of design variables and their mixed (discrete/continuous) nature has rendered the structure optimization a perfect benchmark for large scale search algorithms. The emergence of Evolutionary Algorithms, with their ability to simultaneously examine numerous areas of vast, multimodal search spaces resulted in more reliable and noticeably less expensive designs.

After this introduction, the present paper is organized in the following order: Section 2 is devoted to the general description of the Genetic Algorithm and Nelder Mead Simplex. In Section 3 the hybrid mechanism is fully described and some modifications on the compatibility and efficiency of the algorithm are proposed. Numerical examples and the application of the algorithm are presented in Section 4. Concluding remarks form the content of Section 5.

# 2. GENETIC ALGORITHM AND NELDER-MEAD SIMPLEX

## 2.1. Genetic Algorithm

The genetic algorithm (GA) is a search heuristic that mimic the genetic processes of biological organisms. This heuristic is widely used to explore useful solutions to optimization and search problems. Genetic algorithms belong to the larger class of

evolutionary algorithms (EA), and uses *stochastic direct* optimization methods. The term stochastic indicates that GA uses random operators, and therefore may result in a different (set of) solutions each time they are run. They are also direct, which means they only work with the value of the objective function itself and not its derivatives. GA makes a powerful, universal tool that can solve almost any type of optimization problems including constrained and unconstrained, single- and multi-variable, linear and nonlinear problems with continuous or discrete variables.

They take an initial population of artificial chromosomes and let them evolve toward optimal solution(s) according to the principles of "natural selection" and "survival of the fittest". The evolution usually starts from a population of randomly generated individuals and happens in generations. In each generation, the fitness (measure of desirability) of every individual in the population is evaluated. Multiple individuals are then stochastically selected from the current population (based on their fitness), and modified using such genetic operators as recombination (crossover) and mutation to form a new population. The new population is then used as the "current population" in the next iteration of the algorithm. The average fitness of the new generation is expected to be higher than that of the previous generation, as the best individuals of the last generation have been given higher chances for breeding. Commonly, the algorithm terminates when either a maximum number of generations has been produced, or a satisfactory fitness level has been reached for the population.

Finding the optimal solution of complex high dimensional and multimodal problems often requires very expensive fitness function evaluations. As a result of this, existence of an efficient optimization algorithm seems to be vital. Working with a bunch of solution vectors in each generation, time efficiency has always been a concern in GAs. Other criticisms are also mentioned about the functionality of GA but the worst among these defects is getting trapped in the local optima of the objective function in which an unrealistic solution regarded as the global minima is achieved.

In order to make GA a more reliable algorithm and in relation with optimization efficiency (better solutions with fewer function evaluation) many mechanisms have been suggested. One is improvements upon the mechanism of the algorithm, such as modification of genetic operators, or the use of niche technique, etc; the other is hybridizing of GA with other optimization methods, especially the ones working in local search spaces such as the Nelder-Mead Simplex, simulated annealing (SA), etc.

## 2.2. Nelder Mead simplex

The Nelder–Mead simplex (NMS) search method is based upon the work of Spendley et al. [7]. A simplex is a geometrical figure consisting in n-dimentions, of (n + 1) points:  $x_0, ..., x_n$ . If any point of a simplex is taken as the origin, then other points define vectorer directions that span the n-dimention vector scace. Through a sequence of elementary geometric transformation (reflection, contraction, expansion and multi-contraction), the initial simplex moves, expands or contrasts. To select appropriate transformation, the method only uses the values of the function to be optimizaed at the vertices of the simplex considered. After each transformation, the current worst vertex is replaced by a better one. The trial movements in

Figure 1 are generated according to these opeartors ( $x_h$ : Represents the vertex where the objective function is the highest and  $x_l$  represents the vertex where the objective function is the lowest).



Figure 1. Simplex procedures (reflection, contraction, expansion and multi-contraction)

Operators in simplex method are defined as follow:  $(\bar{x} = \frac{1}{n} \sum_{i=1}^{n} x_i)$ 

- Reflection:  $x_r = (1 + \alpha)\overline{x} \alpha x_0$
- Expansion:  $x_e = (1 \gamma)\overline{x} + \gamma x_r$
- Contraction:  $x_c = (1 \beta)\overline{x} + \beta x_0$

 $\alpha$ ,  $\beta$  and  $\gamma$  are real parameters that control the operators of the Simplex.

Convergence critertion is a measure of how a solution has moved from iteration (k) to (k + 1). So the algorithm stops whenever [8]:

$$\frac{1}{n} \sum_{i=1}^{n} \left\| x_i^{k+1} - x_i^k \right\|^2 < \varepsilon$$
 (1)

Where  $x^{k+1}$  indicates the vertex replacing  $x^k$  at the iteration (k+1), and  $\varepsilon$  is a small real positive number.

The Nelder-Mead simplex method has been applied in physics [9], crystallography [10],

biology [11], chemistry [12] and health care [13]. Fletcher [14] considers the Nelder-Mead technique as one of the most successful methods that merely compare function values. There have been abundant studies devoted to various modifications of the Nelder-Mead simplex method appearing in the literature, such as Barton and Ivey [15], Nazareth and Tzeng [16], etc. However, the simple unmodified Nelder-Mead version of the simplex algorithm is best suited to our needs to hybridize with GA.

## **3. HYBRID ALGORITHM GA-NMS**

### 3.1. Structure of the GA-NMS

In this section the hybrid algorithm GA-NMS is presented in detail. The main idea behind this algorithm is to combine the advantages of each algorithm in a way to avoid their disadvantages. It should also be pointed out that this hybrid algorithm is coded to deal with highly challenging problems as it will be shown in Section 4, and for regular problems considering all of these specifications are not valuable. First the general algorithm is presented and then some modifications are performed to improve the compatibility and efficiency of the algorithm.

#### 3.1.1. Initialization

In the first step the parameters of the algorithm are initialized, these parameters include the simple GA parameters like rate of the Crossover ( $R_c$ ), Mutation probability ( $P_M$ ), the length of the Chromosomes and the number of population, and the controlling parameters that determines the interference of each algorithm (GA and Simplex). Recommendation on the magnitude of GA parameters is beyond the scope of this research, because they are highly dependent on the type of the problem being studied. However some modifications will be presented to improve their efficiency.

#### 3.1.2. Exploration of the search domain

In this section the GA is ignited by the randomly selected chromosome. The algorithm takes these initial points to make up the first mating pool in which the first generation will be born afterwards. Experiments show that a properly coded GA can explore the search domain to the global solution at the first 10,000~30,000 function evaluations, although this magnitude is highly dependent on the complexity of the problem. We want to ignore the common convergence criteria of the GAs by gradually decreasing the interference of the GA in the solution of the algorithm. According to this consideration, in the first steps, the GA is the only working algorithm that travels around the search domain to fine out the optimum regions and may be stopped if one of these criteria is fulfilled.

## 3.1.3. Exploitation

Here we define a controlling parameter which defines the interference of each sub

algorithm in the solution of the main algorithm. We want to transform the solution from an exploration one to exploitation. Thus by calling the interference rate of the GA and the Simplex as  $I_G$ ,  $I_S$  respectively, and by considering that  $I_G +I_S=0.9$ , we aim to transform the solution. The number of solution points that each algorithm works in this section is equal to the number of population multiplied by  $I_G$  or  $I_S$ . In other words if we set a ranked population, we send 0.1 of the best population from this generation to the next generation without any change. This mechanism which is called *elitism* seems to secure the place of the best chromosomes in this transmission. The other part of population is transformed from GA or Simplex according to the steps in Table 1: (A population of 100 is assumed to exist).

Table 1. Operation of controlling operators					
	Step 1	Step 2	Step 3	Step 4	Step 5
Iteration number	100	200	250	300	400
$l_G$	0.9	0.5	0.4	0.2	0 (Stops)
$l_s$	0	0.4	0.5	0.7	0.9
Genetic population	100	60	50	40	0 (Stops)

Table 1. Operation of controlling operators



Figure 2. Schematic diagram of the GA-NMS algorithm

Figure 2 shows a schematic view of the GA-NMS algorithm. In addition in order to make

better generation and mate some of the good features from the best population, we send the best 10% of the population to the GA in order to make better off-springs.

The operation finally changes to a pure exploitation which stops whenever the Simplex convergence critertion defined by Eq. (1) is satisfied. As we see the number of population that goes under the operation of Simplex gradually increases and on the other hand in a simple Simplex method we need at least n+1 vertices (n determines the number of variables). The number of vertices which can be used and the number of separate Simplex that can be run among these points are optional and they can be adjusted by user since they are mainly dependent on the number of population. Experiments show that if we divide the population in a way that leads to a maximum number of sub groups fitted to Simplex algorithm, better results will be achieved.

- 3.2. Modifications
- 3.2.1. Improving GA operators
- 3.2.1.1. Reproduction

A variety of methods have been introduced for reproduction in GA. Goldberg introduced the application of roulette wheel as the most famous tool for this operator. A very easy and more efficient way to choose parents is to randomly choose 2 chromosomes and compare their fitness. The chromosome with better fitness will be the candidate for being a parent. This method is called Tournament Selection. Selection pressure is easily adjusted by changing the tournament size. If the tournament size is larger, weak individuals have a smaller chance to be selected. This method can reproduce better generation as it explores a theme of elitism.

## 3.2.1.2. Dynamic mutation

A method of dynamic establishment of mutation probability in the genetic algorithm can be useful as we aim at exploration property of the GA. The comparison of the results shows an oscillating generations but very swift steps are performed to reach the region of solution. The mixture of reproduction method explained above and the Dynamic Mutation leads to a great combination hoping to achieve the exploration capacity of the GA. A comparison of the common mutation magnitude and the dynamic values is presented in Figure 3. The dynamic value is obtained by easily dividing the  $P_m$  by the generation number. A big value of  $P_m=0.1$  in comparison of ubiquitous magnitudes was considered that gradually decreases in order to stabilize the generation.

### 3.2.2. Compatibility with constrained problems

A simple GA and the Simplex method were both compatible with unconstrained problems. By defining the application of penalty method we transform the problem into an unconstrained one. As we know a constrained problem is one in which the feasible region is defined by a set of implicit/explicit constraints. As an example to optimize weight of a truss we are searching for the lowest cross section of members and the constraints are stress and displacement of members. Here stress and displacement are indirect constraints

as they are not applied directly on the cross section of the members i.e. the variables of the fitness function.



Figure 3. Comparison between static and dynamic mutations

To shift the problem to a normal form, we easily apply the penalty to the solution which violates the constraints. Thus automatically the fitness of these chromosomes are lowered. Many penalty methods are introduced and suggested for different problems. Here we apply the dynamic penalty as follow:

*Obj.Function*: 
$$f(\vec{x})$$
, *Constrained to*:  $g_i(\vec{x}) \le 0, i = 1, ..., q$  (2)

$$\vec{F(x)} = f(\vec{x}) \pm p(\vec{x}) \tag{3}$$

$$p(\vec{x}) = \alpha(\#iteration)^{\beta} \times \sum_{i=0}^{q} S_i(\vec{x}) \text{ in which } S_i(\vec{x}) = \begin{cases} 0 & g_i(\vec{x}) \le 0\\ |g_i(\vec{x})| & otherwise \end{cases}$$
(4)

The sensitivity of the problem to this penalty method is presented in Eqs. (1- 3) and its parameters have direct impact on the final result and are all beyond the scope of this research. Dynamic penalty enables algorithm to have a smooth rate of reaching to the global optimum. A wise tuning of the penalty parameters will tend the algorithm to find the global optimum and preferably search areas nearer to the boundaries of constraints where the validity of the solution is acceptable and the global point also lay there. In Figure 4 two convergence histories for a typical test case are included, one with high penalty orders and the other one with low penalty orders for coefficient in Eq. (3). Knowing the physics of the constraints enables us to adopt proper values, hence converging to more reliable results.



Figure 4. Comparison between the convegence history of penalty methods

### 4. NUMERICAL EXAMPLES

## 4.1. Introduction to structure optimization and dome and large-scale structures

To demonstrate the efficiency of this algorithm, weight minimization of structures is chosen as the test cases. These numerical examples are among those chosen in the literature to show the robustness of different algorithms. The nonlinearity of the problem, number of variables and constraints are the factors that determine the complexity of such problems. The test case problems are chosen in a variety of variables, ranging from 29~200 which are holding a number of constraints up to 3500. In all cases, the presented algorithm overcomes its rivals. For each test case the reported results in the literature are also provided.

From the analytical point of view the weight minimization problem of a trusses structure can be put in the following form:

Fitness function:

$$W(\vec{A}) = \rho g \sum_{j=1}^{N_l} A_j l_j \tag{5}$$

Constrained to:

$$u_{(x,y,z),k}^{l} \leq u_{(x,y,z),k,ilc} \leq u_{(x,y,z),k}^{u} \begin{cases} j = 1, N_{e} \\ k = 1, N_{n} & \& & \sigma_{j}^{l} \leq \sigma_{j,ilc} \leq \sigma_{j}^{u} & \& & A_{j}^{l} \leq A_{j} \leq A_{j}^{u} \\ ilc = 1, N_{c} \end{cases}$$
(6)

Where:

- $N_{e}$  is the number of elements in the structure;
- $N_n$  is the total number of nodes;
- $N_c$  is the number of independent loading conditions acting on the structure;

- g is the gravity acceleration value  $(9.81 \text{ m/s}^2)$ ;
- $\rho$  is the material density;
- $\vec{A} = A(A_1, A_1, A_1, ..., A_N)$  is the vector of the cross sectional areas values;  $A_j$  is the area of the *jth* element of the structure.  $A_j^l$  and  $A_j^u$  are lower and upper bounds of  $A_j$ ;
- $l_i$  is the length of *jth* element, which is:

$$l_{j} = \sqrt{(x_{j1} - x_{j2})^{2} + (y_{j1} - y_{j2})^{2} + (z_{j1} - z_{j2})^{2}}$$
(7)

Where  $x_{j1,2}, y_{j1,2}, z_{j1,2}$  are the nodal coordinates of the *jth* element.

 $u_{(x,y,z),k,ilc}$  are the displacements of the *kth* node in the directions *x*, *y*, *z* with the lower and upper limits  $u_{(x,y,z),k}^{l}$  and  $u_{(x,y,z),k}^{u}$  in corresponding of the *ilcth* loading condition;

•  $\sigma_{j,ilc}$  is the stress on the *jth* element, with the lower and upper limits  $\sigma_k^l$  and  $\sigma_k^u$ , in correspondence of the *ilcth* loading condition.

#### 4.2. Large-scale truss optimization

Three numerical examples chosen from size optimum design of large-scale truss structures are studied to test and verify efficiency of the GA-NMS algorithm, as well as to illustrate its applicability for optimum design of practical structures. The studied trusses consist of a 200 bar truss and the 942 bar truss. These benchmark problems are also studied by different researchers including Farshi and Ziazi[17], Venkaya et al [18], Adeli and Cheng [19], Erbatur and Hasancebi [20]. The general features of these trusses are listed in Table 2.

Property / type of the truss	200-bar truss	942-bar truss
Number of variables	29-200	59
Minimum cross section	$0.1 in^2$	$1 in^2$
Material density	$0.283 \ \frac{1b}{in^2}$	0.1 $\frac{1b}{in^2}$
Modulus of elasticity	30000 ksi	10000 ksi

Table 2. General properties of large-scale trusses

# 4.2.1. A 200-bar truss

Due to the symmetry of the structure, variables linking are adopted and the areas are grouped into twenty-nine groups and 29 optimization variables are considered according to Figure 5. The following three independent loading conditions are imposed on the structure:

(a) 1000 lbf in positive x-direction at nodes 1,6,15,20,29,34,43,48,57,62 and 71;

(b) 10000 lbf in negative y-direction at nodes 1, 2, 3,4,5,6,8,10,

12, 14, 15, 16, 17, 18, 19, 20, 22, 24, 26, 28, 29, 30, 31, 32, 33, 34, 36, 38, 40, 42, 43, 44, 45, 46, 47, 48, 50, 52, 54, 56, 58, 59, 60, 61, 62, 64, 66, 68, 70, 71, 72, 73, 74, 75;

(c) The loading conditions (a) and (b) acting together; Table 3 presents the optimal results for this case study.



Figure 5. Two hundred-bar truss scheme

Variables	Farshi and Ziazi [17]	This research
$A_1(in.^2)$	0.147	0.147
$A_2$	0.945	0.935
$A_3$	0.1	0.1051
$A_4$	0.1	0.1062
$A_5$	1.9451	1.9445
$A_6$	0.2969	0.2982
$A_7$	0.1	0.1
$A_8$	3.1062	3.1064
$A_9$	0.1	0.1092
$A_{10}$	4.1052	4.1089
A <sub>11</sub>	0.4039	0.4031
A <sub>12</sub>	0.1934	0.1957
A <sub>13</sub>	5.4289	5.4201
A <sub>14</sub>	0.1	0.1
A <sub>15</sub>	6.4289	6.4252
A <sub>16</sub>	0.5745	0.5776
A <sub>17</sub>	0.1339	0.1336
A <sub>18</sub>	7.9737	7.9782
A <sub>19</sub>	0.1	0.1
$A_{20}$	8.9737	8.9637
A <sub>21</sub>	0.7053	0.7002
A <sub>22</sub>	0.4215	0.4258
A <sub>23</sub>	10.8675	10.8594
A <sub>24</sub>	0.1	0.1
A <sub>25</sub>	11.8674	11.8636
A <sub>26</sub>	1.0349	1.029
A <sub>27</sub>	6.6849	6.6801
$A_{28}$	10.8101	10.8158
A <sub>29</sub>	13.8379	13.829
Weight (lb)	25456.57	25449.27

Table 3. Optimal designs of the 200-bar truss

4.2.2. The 200 bar space truss structure with five independent loading conditions

The planar 200-bar truss structure shown in Figure 5 can be optimized also with 200 design variables by assigning a sizing variable to the cross-sectional of each element. The structure is designed to carry five independent loading conditions. Besides the three load cases listed in Section 4.2.1, the structure is also loaded by

(d) 1000 lbf acting in the negative x-direction at node points 5,14,19,28,33,42,47,56,61,70 and 75;

(e) Loading conditions (b)–described in Section 4.2.1.1 – and (d) acting together.

The optimization includes 3500 non-linear constraints on nodal displacements and member stresses. The displacements of all free nodes in both directions x and y must be less than  $\pm 0.5$  in. The allowable stress (the same in tension and compression) is 30,000 psi. The lower bound of cross-sectional areas is  $0.1 \text{ in}^2$ .

Table 4 shows the results and a comparison of the available results in the literature.

Venkaya This Venkaya This Venkaya This Venkaya This Venkaya This Variables Variables Variables Variables Variables Work et al[18] Work Work et al[18] et a**l**[18] Work et al[18] Work et al[18] 5 7 37 1 4067 0.233 5.7218 2 558 2 5469 0.1 A 1 348 A40 0.1 Aß A120 A159 0.21 1.313 4.798 4.849 1.988 2.0443 0.237 0.2525 14.9598  $A_2$ 1.1623 Aas A 82 A121 A160 14.981 3 402 3 4 1 8 1 1.85 1.8053 0.201 0.1542 10 649 1.175 10 997 1.2812 As A 44 AB A122 Aia 1.771 1.73 0.127 0.1 7.22 0.7787 1.251 Aδ A<sub>45</sub> A 84 7.4849 0.966 A162 1.1015 A123 4.318 4.2611 0.984 9.8 0.173 0.1328 1.0679 0.991 1.1342 9.7349 AT A46 As A 124 AIB 0.7747 As 1.497 1.5895 Aq 0.971 0.797 0.7845 7.822 7.6835 0.116 0.3897 A 85 A125 A170 0 742 0 4548 0 749 0.6885 5.6498 0.116 0.1 0.816 Ag A48 A 87 5 626 A132 0.4938 Am 0.782 0.8027 3.346 0.116 0.634 0.4905 A10 Aqq 3.162 Agq 0.1 A138 A112 0.816 0.789 0.116 1.156 0.9263 0.2195 0.491 0.7661 0.634 0.3957 0.703 0.6152  $A_{ll}$ Aso Ags A 134 A173 0.491 A12 0.116 0.3031 A51 0.333 0.3364 A 96 0.652 0.512 0.6612 A 178 6.713 6.874 A135 A13 0.377 0.2921 A 58 0.333 0.1501 Ag 0.318 0.4281 A 140 7.285 7.4151 A179 0.713 0.8576 0.587 0.377 0.387 0.208 0.36 6.688 6.5888 0.4113 4.281 4.1989 A14 A 50 A102 Aja A180 5.662 5.605 AIS 0.435 0.3559 A 64 0.533 0.2409 A 142 2.835 2.9147 16.104 16.187 AIR A 181 A16 4 575 4.713 Ass 0.519 0.55 2 1 5 1 2.2602 11.752 11.5395 A182 1.309 1 3174 A 104 Azes 0.538 1.8351  $A_{I7}$ 0.5823 A 66 1.95 A105 8.288 8.4646 A 144 1.049 0.8999 A183 1.317 1.3123 1.895 1.9992 5.326 0.884 1.011 0.921 10.95 Ast 5.3958 0.567 11.2053 A18 A100 A 195 A 189 2.483 2.5937 0.813 0.832 0.984 1.0772 8.969 8.9478 5.073 4.8779 A19 A 68 A107 A 146 A 191 A20 0.75 0.571 0.954 1.084 6.77 6.6058 2.495 2.439 3.243 3.3121 An A153 A 192 A108 0 784 4 4 9 5 4 7932 1 7 2 5 9 1.024 1.0722 8 983 9.017 Agi 0.4193  $A_{70}$ A 115 1.687 A 154 A 195 2.278 2.5342 1.391 1.2508 0.605 0.3831 5.695 5.7271 20.687 20.7437 A-22 An Ans Ain Aigs 3.8349 9,594 1.294 1.2493 0.343 0.1197 6.274 3.932 9,6298 AB A78 A 119 6.4366 A158 Ajg Weight(lb) 30849 31020

 Table 4. Optimal designs of the 200-bar truss with five loading conditions (Areas of elements on one side of symmetrical line)

### 4.2.3. A 942-bar truss

The 26-story tower space truss containing 942 elements and 244 nodes is considered in this section. 59 design variables are used regarding the symmetry of the structure. Figure 6 shows the geometry and the 59 element groups. The members are subjected to the stress limits of  $\pm 25$  ksi ( $\pm 172.375$  MPa) and the four nodes of the top level in the x, y and z directions are subjected to the displacement limits of  $\pm 15.0$ in ( $\pm 38.10$ cm). The allowable cross-sectional areas in this example are selected from 0.1 to 200in<sup>2</sup>. The loading on the structure consists of:

(1) The vertical load at each node in the first section (by section we mean a storey of the structure) is equal to 3 kips. (2) The vertical load at each node in the second section is equal to 6 kips. (3) The vertical load at each node in the third section is equal to 9 kips. (4) The horizontal load at each node on the right side in the x direction is equal to 1 kips. (5) The horizontal load at each node on the left side in the x direction is equal to 1.5 kips. (6) The horizontal load at each node on the front side in the y direction is equal to 1 kips. (7) The horizontal load at each node on the back side in the y direction is equal to 1 kips.

Table 5 lists the available studies on this case and the results for this algorithm. It is obvious that like other cases the results from this algorithm outperforms its rivals.

Variables	Adeli and Cheng [19]	Erbatur and Hasancebi [20]	This research
$A_1$ (in <sup>2</sup> .)	Not Available	1.00	2.7859
A <sub>2</sub>	Not Available	1.00	1.3572
A <sub>3</sub>	Not Available	3.00	5.0362
A <sub>4</sub>	Not Available	1.00	2.2398
As	Not Available	1.00	1.2226
A <sub>e</sub>	Not Available	17.00	14.9575
A <sub>7</sub>	Not Available	3.00	2.9568
A <sub>a</sub>	Not Available	7.00	10.9038
A <sub>9</sub>	Not Available	20.00	14.4177
$A_{16}$	Not Available	1.00	3.7090
A <sub>11</sub>	Not Available	8.00	5.7076
A <sub>12</sub>	Not Available	7.00	4.9264
A <sub>12</sub>	Not Available	19.00	14.1751
A <sub>14</sub>	Not Available	2.00	1.9043
A <sub>15</sub>	Not Available	5.00	2.8101
A <sub>16</sub>	Not Available	1.00	1.0000
A <sub>17</sub>	Not Available	22.00	18.8070
A <sub>15</sub>	Not Available	3.00	2.6151
A <sub>19</sub>	Not Available	9.00	12.5328
A <sub>26</sub>	Not Available	1.00	1.1314
A <sub>21</sub>	Not Available	34.00	30.5122
A <sub>22</sub>	Not Available	3.00	3.3460
A <sub>23</sub>	Not Available	19.00	17.0450
A <sub>24</sub>	Not Available	27.00	18.0785
A <sub>28</sub>	Not Available	42.00	39.2717
Aze	Not Available	1.00	2.6062
A <sub>27</sub>	Not Available	12.00	9.8303
A <sub>28</sub>	Not Available	16.00	13.1126
A <sub>29</sub>	Not Available	19.00	13.6897
A <sub>30</sub>	Not Available	14.00	16.9776

Table 5. Optimal design of 942-bar tower space truss

Variables	Adeli and Cheng [19]	Erbatur and Hasancebi [20]	This research
A <sub>31</sub>	Not Available	42.00	37.6006
A <sub>32</sub>	Not Available	4.00	3.0602
A <sub>33</sub>	Not Available	4.00	5.5106
A <sub>34</sub>	Not Available	4.00	1.8014
A <sub>35</sub>	Not Available	1.00	1.1568
A <sub>36</sub>	Not Available	1.00	1.2423
A <sub>37</sub>	Not Available	62.00	62.7741
Ass	Not Available	3.00	3.3276
A <sub>39</sub>	Not Available	2.00	4.2369
A <sub>46</sub>	Not Available	4.00	1.7202
A <sub>41</sub>	Not Available	1.00	1.0148
A <sub>42</sub>	Not Available	2.00	5.6428
A <sub>43</sub>	Not Available	77.00	78.0094
A <sub>44</sub>	Not Available	3.00	3.2206
A <sub>45</sub>	Not Available	2.00	3.5934
A <sub>46</sub>	Not Available	3.00	4.7668
A <sub>47</sub>	Not Available	2.00	1.1531
A <sub>48</sub>	Not Available	3.00	2.1698
A <sub>49</sub>	Not Available	100.00	99.6406
Aac	Not Available	4.00	4.1469
A <sub>51</sub>	Not Available	1.00	2.1600
A <sub>52</sub>	Not Available	4.00	4.1499
A <sub>33</sub>	Not Available	6.00	11.2070
A <sub>34</sub>	Not Available	3.00	11.0904
Aas	Not Available	49.00	35.9499
Aza	Not Available	1.00	2.1937
A <sub>57</sub>	Not Available	62.00	66.1705
A <sub>SE</sub>	Not Available	1.00	3.3402
A <sub>ag</sub>	Not Available	3	4.0525
Weight (lb)	170000	143436	142295.75

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Figure 6. A 26-story tower space truss

## **5. CONCLUSION**

In this paper a hybrid algorithm is presented for finding the global minimum of the objective function for both constrained and unconstrained problems. A controlled mixture of main characteristics of two different algorithms (exploration of Genetic algorithm and exploitation of Nelder-Mead simplex) led to a new algorithm which has promising results in the challenging field of structural optimization. Moreover some modifications are implemented on the GA operators to improve its convergence rate and reaching better results. Optimization costs in terms of CPU usage or convergence time decreased since simple-coded algorithm like Nelder-Mead simplex were used in hybridizing with modified GA.

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